

Kiama: Domain-Specific Languages for Language Implementation in Scala

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Language processing problems

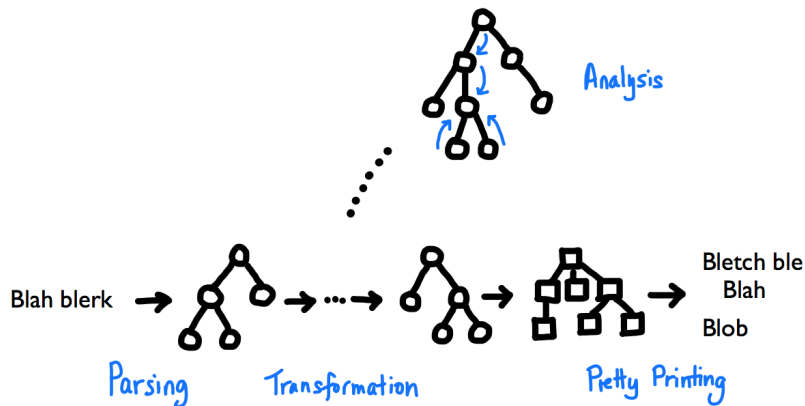


Figure : The big picture

A collection of Scala-based *internal domain-specific languages* for solving problems in the domain of language processing.

<i>Problem</i>	<i>Domain-specific language</i>
Text to structure	Context-free grammars
Structure representation	Algebraic data types
Transformation	Term rewriting
Analysis	Attribute grammars
Structure to text	Pretty printing

Lambda calculus REPL

```
> 1 + 2
3 : Int

> (\x : Int . x + 1) 4
5 : Int

> let i : Int = ((\x : Int . x + 1) 4) in i + i
10 : Int

> \x : Int . x + y
1.16: y: unknown variable

> \x : Int . x 4
1.12: x: non-function applied

> (\f : Int -> Int . f 4) 5
1.25: 5: expected Int -> Int, got Int
```

Abstract Syntax Trees

```
let i : Int = ((\x : Int . x + 1) 4) in i + i
```

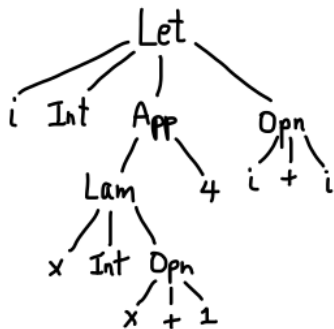


Figure : The data is naturally represented by a hierarchical data structure

Evaluation: beta reduction rule

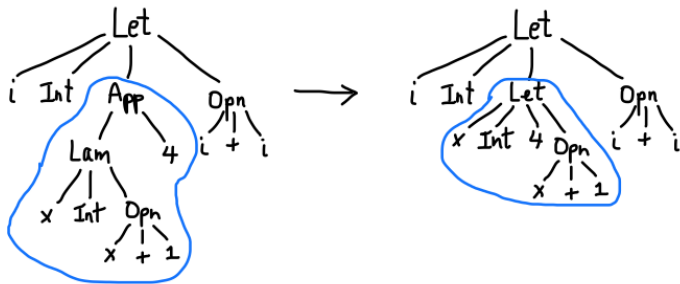


Figure : Replace an application by a substitution into the body

```
val beta =  
  rule {  
    case App (Lam (x, t, e1), e2) =>  
      Let (x, t, e2, e1)  
  }
```

Evaluation: substitution rule

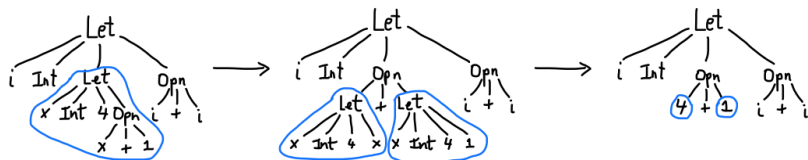


Figure : Push substitutions into operations, replace variables

```
val substitution =  
  rule {  
    case Let (x, t, e, Opn (l, op, r)) =>  
      Opn (Let (x, t, e, l), op, Let (x, t, e, r))  
    case Let (x, _, e, v @ Var (y))    =>  
      if (x == y) e else v  
    case Let (_, _, _, n : Num)       => n  
    ...  
  }
```

Choosing between rewrite strategies

```
val onestep : Strategy =  
  beta <+ arithop <+ substitution
```

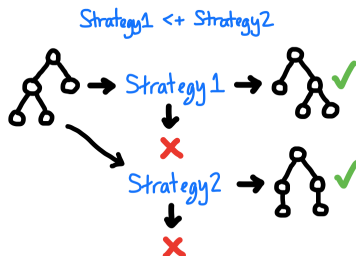


Figure : A choice succeeds iff one of its alternatives succeeds

Generic traversal: some

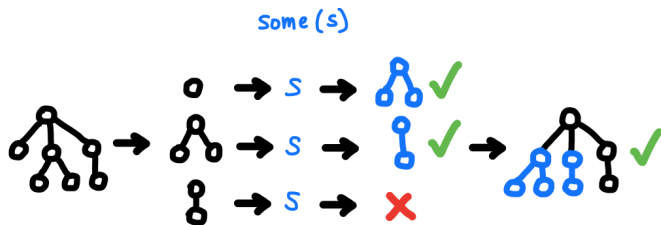


Figure : some (s) succeeds if s succeeds at least one of the children

Building up a traversal

- ▶ At the root

```
onestep
```

- ▶ At one or more children of the root

```
some (onestep)
```

- ▶ At one or more children of the root, or if that fails, at the root

```
some (onestep) <+ onestep
```

- ▶ At the topmost nodes at which it succeeds

```
lazy val inner = some (inner) <+ onestep
```

Type checking: attributes



Figure : Attributes are memoised functions of nodes

```
val envOf : Exp => List[(String,Type)] =  
  attr {  
    ... cases go here ...  
  }  
  
val typeOf : Exp => Type =  
  attr {  
    ... cases go here ...  
  }
```

Environments

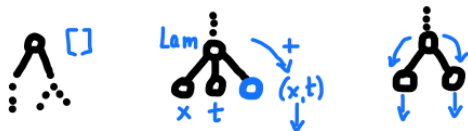


Figure : Propagate downward, gather bindings at lambdas and lets

```
case n if n.isRoot                => List ()
case n =>
  n.parent match {
    case p @ Lam (x, t, _)         => (x,t) :: envOf (p)
    case p @ Let (x, t, _, _)     => (x,t) :: envOf (p)
    case p : Exp                  => envOf (p)
  }
```

Type checking: numbers, lambdas and lets

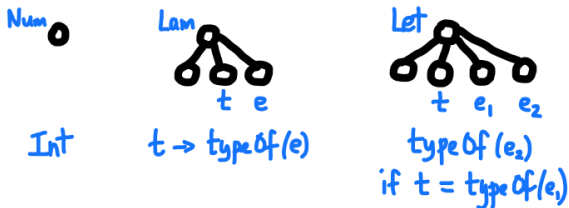


Figure : Lambdas and lets get types from their children

```
case Num (_) => IntType ()
case Lam (_, t, e) => FunType (t, typeOf (e))
case Let (_, t, e1, e2) => if (checkType (e1, t))
                             typeOf (e2)
                             else
                               UnknownType ()
```

Type checking: application

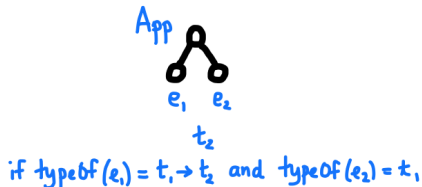


Figure : Applications get types from the applied functions

```
case App (e1, e2) =>
  typeof (e1) match {
    case FunType (t1, t2) =>
      if (checkType (e2, t1)) t2
      else UnknownType ()
    case _ =>
      message (e1, e1 + ": non-function applied")
      UnknownType ()
  }
```

Type checking: variables

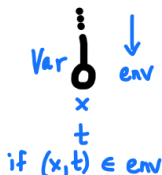


Figure : Variables get their types from the environment

```
case e @ Var (x) =>
  envOf (e).collectFirst {
    case (y, t) if x == y => t
  }.getOrElse {
    message (e, x + ": unknown variable")
    UnknownType ()
  }
```

Why should I care?

- ▶ Lambda calculus is fun!
 - ▶ bitbucket.org/inkytonik/lambdajam13
- ▶ Many problems take this form:
 - ▶ extracting information from structured data
 - ▶ checking that structured data meets constraints
 - ▶ translating structured data into other forms
 - ▶ using analysis results to guide transformation
- ▶ Abstracting away from details of traversal and storage
 - ▶ focusses your attention on the transformations and the analyses,
 - ▶ simplifies what you have to write, and
 - ▶ makes the computation more independent of the structure.

Conclusion

- ▶ Kiama abstracts language processing problems by tree rewriting and tree decoration.
- ▶ Details of traversal and storage are almost entirely implicit.
- ▶ The focus is on the transformations and analyses.
- ▶ Scala enables these DSLs to be implemented easily and they integrate naturally with other code.
- ▶ Wait there's more!
 - ▶ rewriting collections
 - ▶ pretty-printing domain-specific language
 - ▶ frameworks for compilers, REPLs and testing
 - ▶ augmentation of Scala parser combinator library
 - ▶ profiling of rewriting and attribution

More Information

- ▶ Kiama project: kiama.googlecode.com
 - ▶ Papers and talks: code.google.com/p/kiama/wiki/Research
- ▶ Other projects:
 - ▶ Using macros to get names: bitbucket.org/inkytonik/dsinfo
 - ▶ Domain-specific profiling: bitbucket.org/inkytonik/dsprofile
 - ▶ sbt plugin for Rats! parser generator: sbt-rats.googlecode.com
 - ▶ Scala worksheet for Sublime Text 3:
bitbucket.org/inkytonik/scalaworksheet
- ▶ Related work:
 - ▶ Stratego: strategoxt.org
 - ▶ JastAdd: jastadd.org
 - ▶ Swierstra, S., and Chitil, O. *Linear, bounded, functional pretty printing*. *Journal of Functional Programming* 19, 1 (2008), 1–16

Workshop

- ▶ Lambda calculus
 - ▶ New feature:
 - ▶ matching by cases
 - ▶ New processing:
 - ▶ different evaluation strategies
 - ▶ profiling demonstration
- ▶ Kiama
 - ▶ Implementation of rewriting and attribution DSLs
 - ▶ How do we use Scala macros?
 - ▶ Pretty-printing DSL
 - ▶ How does profiling work?